

Generating New Prime Numbers via Symmetric Triplets: 20 New 300-Digit Primes in 40 Seconds

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Abstract:

This paper presents a novel method for discovering large prime numbers using symmetric triplets, based on an extension of the Goldbach hypothesis. In a previous paper (DOI: 10.5281/zenodo.15198172, <https://ai.vixra.org/abs/2504.0029>), we demonstrated that every positive integer is surrounded by a symmetric pair of prime numbers. In this paper, we examine a specific case of that symmetry—where the center of symmetry is itself a prime number. Thus, prime numbers can be grouped into “triplets”: two primes equidistant from a third one. The key idea is to analyze the most common differences between primes and to scale these differences based on the distribution density, enabling efficient discovery of large primes with far fewer checks compared to brute-force methods.

1. Introduction

Prime numbers are fundamental in number theory and essential to modern cryptography. Generating large primes is crucial for RSA and other encryption systems, but traditional methods are slow and computationally intensive. This work presents a new approach that significantly reduces the search space using symmetry and statistical inference.

2. Theoretical Basis

The method builds on an extension of Goldbach’s conjecture, analyzing the symmetry of primes around a number n . It assumes that for any integer n , including very large ones, there exists a symmetric pair of primes p_1 and p_2 such that:

$$p_1 + p_2 = 2n \rightarrow n = (p_1 + p_2)/2 \rightarrow p = 2n - s$$

where s is a known prime. This allows the generation of candidate primes through symmetric transformation around n . But this is the general principle that applies to every positive integer. Here, we focus on applying it specifically to known prime numbers, with the goal of finding the “symmetric” ones around them. This reduces the number of required checks.

3. Algorithm and Implementation

The algorithm consists of:

1. Selecting a large base prime n (e.g., verified using FactorDB).

2. Using the most frequent prime gaps up to 1 million: 6, 12, 30, 60, 90, 210, 420, 630, 840, 1050.

3. Scaling each gap to match the size of n using the logarithmic approximation based on Riemann:

$$\text{coefficient} = \text{li}(n) / \text{li}(10^6)$$

4. For each scaled gap d , computing values in a 1% extended interval $d \pm 1\%$, and checking up to 200 million candidates (100 million in each direction). These are not full interval scans but samples spread uniformly across the range. These values are chosen because the scaling, even when based on the Riemann Hypothesis, cannot predict a new prime number with absolute precision. There is an inevitable element of inaccuracy, but as demonstrated in this work, a $\pm 1\%$ interval is sufficient to discover a new prime. Moreover, it is not necessary to check the entire interval—only 200 million values evenly distributed within it.

5. Checking whether $p1 = n - d$ and $p2 = n + d$ are both prime.
6. Running the process in parallel and stopping after discovering a predefined number of new primes.

4. Experiment and Results

Using a 300-digit base prime n , and 10 scaled gaps with 100 million values per gap, we obtained:

- ~2 billion planned checks
- Runtime: 40 seconds (on an 8-core Google Cloud server), stopped early after 10 prime pairs were found. That is, in reality, 2 billion checks were not actually performed.
- New primes discovered: 20 (in 10 symmetric pairs)
- Confirmed using `gmpy2.is_prime` and FactorDB

5. Discussion

Results show that even with modest hardware, proper selection of symmetric candidates and statistical scaling can yield large new primes efficiently. The method is suitable for practical cryptography using primes with 500–1000 digits. Unlike traditional probabilistic methods, this approach offers guaranteed prime numbers in seconds. Although initial screening uses Miller-Rabin, the final numbers were confirmed as primes by independent databases. This opens the door for generating deterministic cryptographic keys in real applications.

6. Scaling Up

The idea of "symmetric primes" would undoubtedly lead to faster discovery of prime numbers across any range when compared to brute-force methods. However, in the higher ranges (with millions of digits), the method becomes impractical due to the immense computational cost of verifying each candidate. Even with a preliminary Miller-Rabin test, checking candidates with millions of digits would take an extremely long time. Additionally, the number of checks must be scaled. For instance, while a 100-million-value interval may be sufficient around each candidate at the 300-digit level, applying a density-based scaling factor means that around 130 trillion checks would be needed at the 20-million-digit level. This alone is not unreasonable, but the main issue is the time required for each primality test. At such digit lengths, each test would take a prohibitive amount of time. Even with supercomputers, it would require thousands or millions of years. Again, this is still faster than brute-force, but in practice, it remains infeasible. Therefore, the method is most effective for practical cryptographic purposes, where primes with 500–1000 digits are used. In that range, truly prime numbers can be generated within seconds.

7. Empirical Results

The following 20 new 300+ digit prime numbers were discovered in 40 seconds on an 8-core cloud instance. All have been submitted and verified as prime by FactorDB:

Base prime number:

p =
1088886711584158931203436278137248555082792661571171969770354590348413
2043015623721673322604207052785909750803572004491817501408260167607773
3788073677923712749770404015724196960297257347124688307337323887078012
4678159746785949926624551330167079023022486903438233545719222702110426
85572248211792573079822078133

Found symmetric primes:

p1 =
1088886593975024039548478314496162359937439424311497432285604497528431
9294130471251521233782691042001304666190109216203204066004803628252097
5979738793566495363516152344305019744769823373493869907585000745202836
5445129424952216811968543302047780043616010279856600792813762430883507
56406901084143384315836829877

p2 =
1088886829193293822858394241778334750228145898830846507255104683168394
4791900776191825411425723063570514835417034792780430936811716706963449
1596408562280930136024655687143374175824691320755506707089647028953188
3911190068619683041280559358286378002428963527019866298624682973337346
14737595339441761843807326389

.....

p1 =

1088886594854807209731270458915260948086653874807022043579676371684280
0967112074102214126066399330673234122912449871473703720702357035695333
0265626545108589196030744184436387946676153111963470701908935687432728
4621216165583952898450964051708493119927018124526522242448927263733022
55073803096009294723179155637

p2 =

1088886828313510652675602097359236162078931448335321895961032809012546
3118919173341132519142014774898585378694694137509931282114163299520213
7310520810738836303510063847012005973918361582285905912765712086723296
4735103327987946954798138608625664926117955682349944848989518140487831
16070693327575851436465000629

.....

p1 =

1088886595439404822179428129429693518562890186799172278030008650443481
4117103791987315999817543934421795689536945902168115945251333828857172
1464801121834843664765256100771424580689112065729024592757753023578718
1620265431524423923886658706935053448324127449820302995380499641117547
87398498743196187191373009077

p2 =

1088886827728913040227444426844803591602695136343171661510700530253344
9968927455456030645390870171150023812070198106815519057565186506358374
6111346234012581834775551930676969339905402628520352021916894750577306
7736054062047475929362443953399104597720846357056164096057945763103305
83745997680388958968271147189

.....

p1 =

1088886594850310025629495520090650133265908894131465204709544148479809
3223582699086756524794801225356299031828218858683360062239302727104497
1938516605995534184997156993549144444137171065066382593313951347588642
1543929296946513440289291290489142053804753168287410201291929708699244
29146421455985417379521826997

p2 =

1088886828318007836777377036183846976899676429010878734831165032217017
0862448548356590120413612880215520469778925150300274940577217608111049
5637630749851891314543651037899249476457343629182994021360696426567382
7812390196625386412959811369845015992240220638589056890146515695521609
41998074967599728780122329269

.....

p1 =

1088886594850021453715808181160946076889406241186064584319733936446925
7262817167143790708576101538700986493868794612639514530792770456457106
6285147202992821577936032008011348976637745349691122154272581014006002
0536767776104700541666165843335991260415197665369321842580888383345112
00483015444359764756601573557

p2 =

1088886828318296408691064375113551033276179081956279355220975244249900
6823214080299555936632312566870833007738349396344120472023749878758440
1291000152854603921604776023437044943956769344558254460402066760150022
8819551717467199311582936816998166785629776141507145248857557020875741
70661480979225381403042582709

.....

p1 =

1088886594556800057344702656073532846981651718192154513275813517854418
5955586982414869702895112791416993727442126644544840709475566835853237
5520648631452434904918399803258202567098233091919866277549698932903180
2287498087825830435292212868911926268481293816227353306598852677045300
58742979987263434952368263349

p2 =

1088886828611517805062169900200964263183933604950189426264895662842407
8130444265028476942313301314154825774165017364438794293340953499362309
2055498724394990594622408228190191353496281602329510337124948841252844
7068821405746069417956889791422231777563679990649113784839592727175553
12401516436321711207275892917

.....

p1 =

1088886593383547204914674470691176877439374916142644501274055129018909
1522704297721114501491409311891055409090948240527209402364332507814030
3740249624379148447460064270174786779131672324509948564847624016046526
8489761116888979322393543082896775850328680094740018622651666654790946
05085013081086529621465305269

p2 =

1088886829784770657492198085583320232726210406999699438266654051677917
2563326949722232143717004793680764092516195768456425600452187827401516
3835897731468277052080743761273607141462842369739428049827023758109498

0866558376682920530855559577437382195716293712136448468786778749429907
66059483342498616538178850997

.....

p1 =

1088886593380809138680738648485185502186708766613422020745150252858775
1135940889028744513115060261802622357301684296956087768275850195508126
7121013693697807758858998929060929090877854979265788057496062169552445
1972852981155593570441565009247859565873320478217389590181174095222369
52742652294300225310709321909

p2 =

1088886829787508723726133907789311607978876556528921918795558927838051
2950090358414602132093353843769197144305459712027547234540670139707420
0455133662149617740681809102387464829716659714983588557178585604603579
7383466512416306282807537651086298480171653328659077501257271308998484
18401844129284920848934834357

.....

p1 =

1088886593967194798923537696976972294199500676051319108043467250323851
2624376221841833747590903452930055850990737872023754689191177364593709
1216031545061082142452845025986281265622555974015630377085115683392353
8482393792298888662552465869702839215917656705045640969204426218752177
53431344960576920220971634869

p2 =

1088886829201123063483334859297524815966084647091024831497241930372975
1461655025601512897617510652641763650616406136959880313625342970621837
6360115810786343357087963005462112654971958720233746237589532090763671
0873925701273011190696636790631318830127317101830826122234019185468676
17713151463008225938672521397

.....

p1 =

1088886593379772001978526492027352539893066694319267208369337430595134
6026119133750799919297625028186679418007590706142935546369930714123565
0431092424213048125739987645594037590165127105728899386485336242968834
6204461128864637972487986956824499070728691711580588849478889702109578
70860477842131227679356162229

p2 =

1088886829788545860428346064247144570272518628823076731171371750101691

8059912113692546725910789077385140083599553302840699456446589621091981
7145054931634377373800820385854356330429387588520477228189311531187190
3151858364707261880761115703509658975316282095295878241959555702111275
00284018581453918480287994037

8. Code Used

The following Python code was used for the purposes of the experiment:

```
import gmpy2

from sympy import li

from time import time, sleep

from multiprocessing import Process, Queue, cpu_count, Value, Lock

from numpy import linspace

# === SETTINGS ===

BASE_PRIME =
gmpy2.mpz("108888671158415893120343627813724855508279266157117196977035459
03484132043015623721673322604207052785909750803572004491817501408260167607
77337880736779237127497704040157241969602972573471246883073373238870780124
67815974678594992662455133016707902302248690343823354571922270211042685572
248211792573079822078133")

BASE_SAMPLE = gmpy2.mpz(10**6)

li_now = li(BASE_PRIME)

li_sample = li(BASE_SAMPLE)

scaling_factor = float(li_now / li_sample)

COMMON_DIFFS = [6, 12, 30, 60, 90, 210, 420, 630, 840, 1050]

scaled_diffs = [int(d * scaling_factor) for d in COMMON_DIFFS]

INTERVAL_PERCENT = 0.01

SAMPLES_PER_DIFF = 100_000_000 # 100 million per difference

print("Safe Scan mode: Scanning 100 million values per gap using parallel cores.")
```

```
# ===WORKER FUNCTIONS ===
```

```
def check_segment(diff_value, start_idx, end_idx, steps, queue, counter, lock):
```

```
    delta = int(diff_value * INTERVAL_PERCENT)
```

```
    start_d = diff_value - delta
```

```
    end_d = diff_value + delta
```

```
    for i in range(start_idx, end_idx):
```

```
        with lock:
```

```
            if counter.value >= 10:
```

```
                return
```

```
            interp_d = int(start_d + i * (end_d - start_d) / steps)
```

```
            p1 = BASE_PRIME - interp_d
```

```
            p2 = BASE_PRIME + interp_d
```

```
            if gmpy2.is_prime(p1) and gmpy2.is_prime(p2):
```

```
                with lock:
```

```
                    if counter.value < 10:
```

```
                        counter.value += 1
```

```
                        queue.put((int(p1), int(BASE_PRIME), int(p2)))
```

```
# === WRITER PROCESS ===
```

```
def writer(queue):
```

```
    with open("found_primes.txt", "a") as f:
```

```
        while True:
```

```
            item = queue.get()
```

```
            if item == "DONE":
```

```
                break
```

```
            p1, p, p2 = item
```

```
f.write(f"p1 = {p1}\np = {p}\np2 = {p2}\n\n")
```

```
# === START ===
```

```
if __name__ == "__main__":
```

```
    start_time = time()
```

```
    cpu_cores = cpu_count()
```

```
    q = Queue()
```

```
    counter = Value('i', 0)
```

```
    lock = Lock()
```

```
    writer_process = Process(target=writer, args=(q,))
```

```
    writer_process.start()
```

```
    for diff_index, d in enumerate(scaled_diffs):
```

```
        print(f"\n → Starting difference #{diff_index + 1} ({d})")
```

```
        segment_size = SAMPLES_PER_DIFF // cpu_cores
```

```
        workers = []
```

```
        for core in range(cpu_cores):
```

```
            start_idx = core * segment_size
```

```
            end_idx = (core + 1) * segment_size if core < cpu_cores - 1 else SAMPLES_PER_DIFF
```

```
            p = Process(target=check_segment, args=(d, start_idx, end_idx,  
SAMPLES_PER_DIFF, q, counter, lock))
```

```
            workers.append(p)
```

```
            p.start()
```

```

    for p in workers:

        p.join()

    with lock:

        if counter.value >= 10:

            print("● 10 pairs found, stopping the search.")

            break

    q.put("DONE")

    writer_process.join()

end_time = time()

print("\n--- RESULTS ---")

print(f"Completed in {end_time - start_time:.2f} seconds ")

```

9. Future Work

1. Testing on more powerful supercomputers to validate scalability to larger ranges.
2. Exploring application of the method to Mersenne numbers, which are easier to verify and may enable breakthroughs at higher levels without massive infrastructure.

10. Conclusion

The presented method enables discovery of large prime numbers using symmetric triplets and scaled prime gaps. It demonstrates practical efficiency and has potential to extend beyond 10^{500} or 10^{1000} .

11. Acknowledgments

The author expresses gratitude to GPT-4 (OpenAI ChatGPT), whose interactive assistance helped resolve many of the technical questions and computations.

12. References:

- Goldbach's conjecture

- Prime number theorem (logarithmic behavior)
- Riemann prime counting function (log integral approximation)
- Miller-Rabin primality test (probabilistic)
- gmpy2.is_prime — fast primality check in Python
- FactorDB for prime number confirmation
- Previous publication: <https://ai.vixra.org/abs/2504.0029>

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